

# A ruler regularity in hexadecimal games

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An important problem in the theory of impartial games is to determine the regularities of their nim-sequences. Subtraction games have periodic nim-sequences and those of octal games are conjectured to be periodic, but the possible regularities of the nim-sequence of a hexadecimal game are unknown. Periodic and arithmetic periodic nim-sequences have been discovered but other patterns also exist. We present an infinite set of hexadecimal games, based on the game **0.2048**, that exhibit a regularity — ruler regularity — not yet reported or codified.

## 1. Introduction

A taking-and-breaking game [Albert et al. 2007; Berlekamp et al. 2001] is an impartial combinatorial game, played with heaps of beans on a table. A move for either player consists of choosing a heap, removing a certain number of beans from the heap, and then possibly splitting the remainder into several heaps; the winner is the player making the last move. For example, both Grundy's Game (choose a heap and split it into two unequal heaps) and Couples-Are-Forever (choose a heap with at least three beans and split it into two) are taking-and-breaking games with very simple rules, however neither has been solved.

We present an overview of the required theory of *impartial* games. The reader can consult the references above for a more in-depth grounding in the theory of, and for more details about, subtraction and octal games.

The *followers* of a game are all those positions which can be reached in one move. The *minimum excluded value* of a set  $S$  is the least nonnegative integer which is not included in  $S$  and is denoted  $\text{mex}(S)$ . The *nim-value* of an impartial game  $G$ , denoted by  $\mathcal{G}(G)$ , is given by  $\mathcal{G}(G) = \text{mex}\{\mathcal{G}(H) \mid H \text{ is a follower of } G\}$ . The values in the set  $\{\mathcal{G}(H) \mid H \text{ is a follower of } G\}$  are called *excluded values* for  $\mathcal{G}(G)$ . An impartial game  $G$  is a previous player win (i.e., the next player has no good move) if and only if  $\mathcal{G}(G) = 0$ . The *exclusive or* (XOR), or *nim-sum*, of two nonnegative integers  $a, b$ , written  $a \oplus b$ , consists of adding their binary

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representations with no carrying; for example,  $3 \oplus 6 = (11 \oplus 110)_{\text{base } 2} = 101_{\text{base } 2} = 5$ . The *disjunctive sum* of games  $G$  and  $H$ , written  $G + H$ , is the game where a player chooses either of the two and plays in that game; the nim-value of a disjunctive sum satisfies  $\mathcal{G}(G + H) = \mathcal{G}(G) \oplus \mathcal{G}(H)$ . Taking-and-breaking games are examples of disjunctive games—choose one heap and play in it. To know how to play these game well, it suffices to know what the nim-values are for individual heaps. For a given game  $G$ , let  $\mathcal{G}(i)$  be the nim-value of  $G$  played with a heap of size  $i$ . We define the nim-sequence for a taking-and-breaking game to be the sequence  $\mathcal{G}(0), \mathcal{G}(1), \mathcal{G}(2), \dots$ . A nim sequence is *periodic* if there exist  $N$  and  $p$  such that  $\mathcal{G}(n + p) = \mathcal{G}(n)$  for all  $n \geq N$ . It is *arithmetic-periodic* if there exist  $N$ ,  $p$  and  $s$  such that  $\mathcal{G}(n + p) = \mathcal{G}(n) + s$  for all  $n \geq N$ , where  $s$  is called the *saltus*.

In a hexadecimal game, after removing beans from a heap the remainder can be split into at most three heaps. The rules for a hexadecimal game are described by a hexadecimal code  $\mathbf{0}.d_1d_2 \dots d_u$  where  $0 \leq d_i \leq 15$ . This is an extension of the octal code used in [Berlekamp et al. 2001]. We use the letters  $A, B, C, D, E, F$  for the numbers 10 through 15 respectively. If  $d_i = 0$  then a player cannot take  $i$  beans away from a heap. If  $d_i = a_32^3 + a_22^2 + a_12^1 + a_02^0$  where  $a_j$  is 0 or 1, a player can remove  $i$  beans from the heap provided he leaves the remainder in exactly  $j$  heaps for some  $j$  with  $a_j = 1$ . If  $0 \leq d_i \leq 7$  for all  $i$  then this is called an *octal* game. This restriction allows a heap to be split into no more than 2 heaps. A subtraction game has  $d_i = 0$  or 3; that is, a player can remove beans but cannot split the heap.

The nim-sequence of a subtraction game with a finite subtraction set is periodic [Albert et al. 2007; Berlekamp et al. 2001]. In [Althöfer and Bültmann 1995], it is shown that the sequence of nim-values of games with small subtraction sets can have long periods. It is conjectured that all octal games also have periodic nim-sequence in unsolved problem 2 of [Guy and Nowakowski 2002]. This appears to be hard. For example, in the game **0.106** Flammenkamp computed the period length as 328226140474 and preperiod length as 465384263797. See [Flammenkamp 2012] for this and other searches for long periods, also see [Berlekamp et al. 2001; Caines et al. 1999; Gangolli and Plambeck 1989]. It is interesting to note that while octal games cannot have arithmetic-periodic nim-sequences [Austin 1976], if a single pass move is allowed then not only do arithmetic-periodic nim-sequences occur, but also nim-sequences composed of a periodic subsequence and an arithmetic-periodic subsequence occur called *sapp* sequences [Horrocks and Nowakowski 2003].

For hexadecimal games, there is an even richer selection of behaviours [Howse and Nowakowski 2004]):

- *Periodic*: **0.B**, period 2 with no preperiod.

- *Arithmetic-periodic*: **0.137F** has nim-sequence  $0, 1, 1, 2, 2, 3, 3, \dots$ , where  $\mathcal{G}(2m-1) = \mathcal{G}(2m) = m$  and  $\mathcal{G}(2m+1) = \mathcal{G}(2m) + 1 = m + 1$  for  $m \geq 1$ . In this case, the saltus is 1 and the period length is 2.
- *Sapp* regularity: **0.205200C** has the nim-sequence consisting of
  - (a) periodic subsequences  $\mathcal{G}(40k+19) = 6$  and  $\mathcal{G}(40k+39) = 14$ , and
  - (b) arithmetic-periodic subsequences

$$\mathcal{G}(40k+j) = \mathcal{G}(40(k-1)+j) + 16, \quad j \neq 19, 39, k \geq 0,$$

with preperiod length of 4, period length of 40.

- In [Howse and Nowakowski 2004], it is noted that **0.123456789** exhibits another type of regularity. Starting with  $n = 0$ , the first fifteen nim-values are  $0, 1, 0, 2, 2, 1, 1, 3, 2, 4, 4, 5, 5, 6, 4$ , and thereafter

$$\mathcal{G}(2m-1) = \mathcal{G}(2m) = m-1,$$

except  $\mathcal{G}(2^k+6) = 2^k - 1$ . The nim-sequence is essentially arithmetic-periodic, but with an infinite number of exceptional values that occur in a geometric fashion.

In this paper, we report another new regularity: *ruler-regularity*. This concept is taken from the markings of a foot ruler. (See [Berlekamp et al. 2003, page 470, Figure 7], “The  $\mathcal{G}$ -values for the RULER game”, except for us the ruler needs to be rotated by 45 degrees and then flipped!)

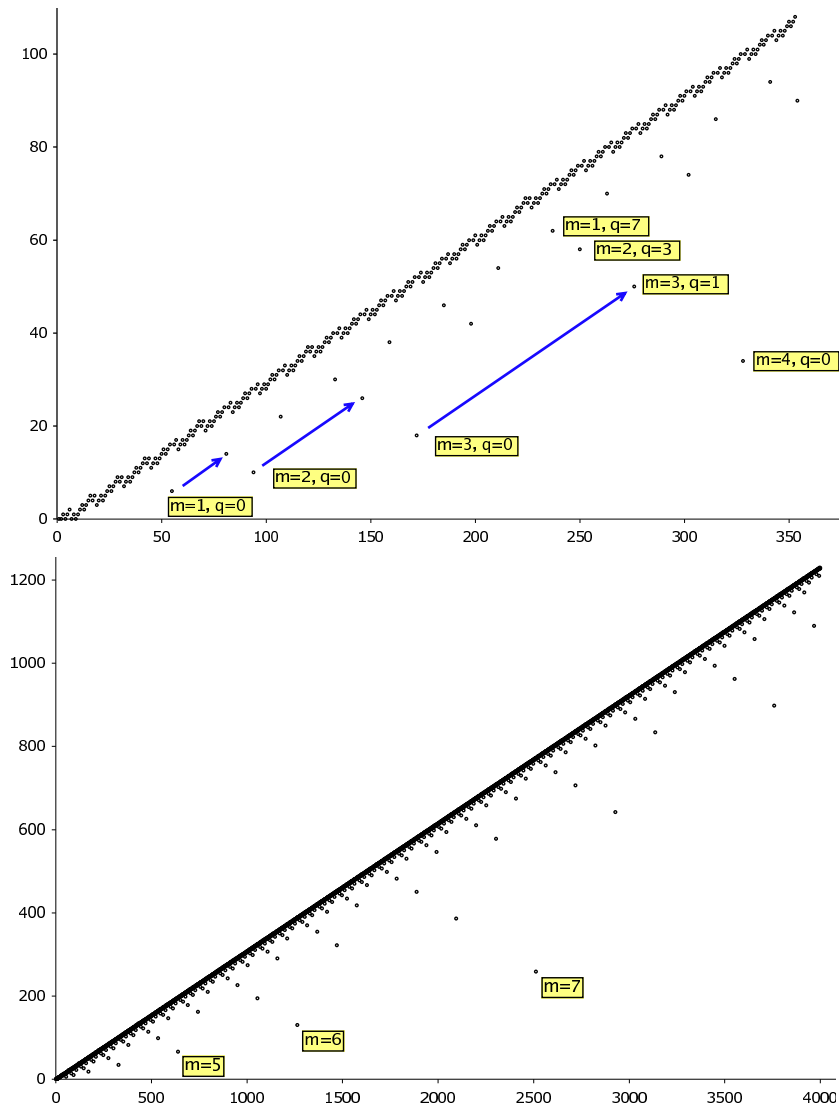
**Definition 1.** A sequence is *ruler regular* if there are positive integers  $N, p, s, r$  and a finite set of integers  $K$  such that  $\mathcal{G}(n+p) = \mathcal{G}(n) + s$  for all  $n \geq N$  except

$$n = r((q+1)2^{m+1} + 2^m + 1) + k, \quad k \in K, q, m \geq 0,$$

when

$$\mathcal{G}(r((q+1)2^{m+1} + 2^m + 1) + k) = \mathcal{G}(r(q2^{m+1} + 2^m + 1) + k) + 2^{m+3}.$$

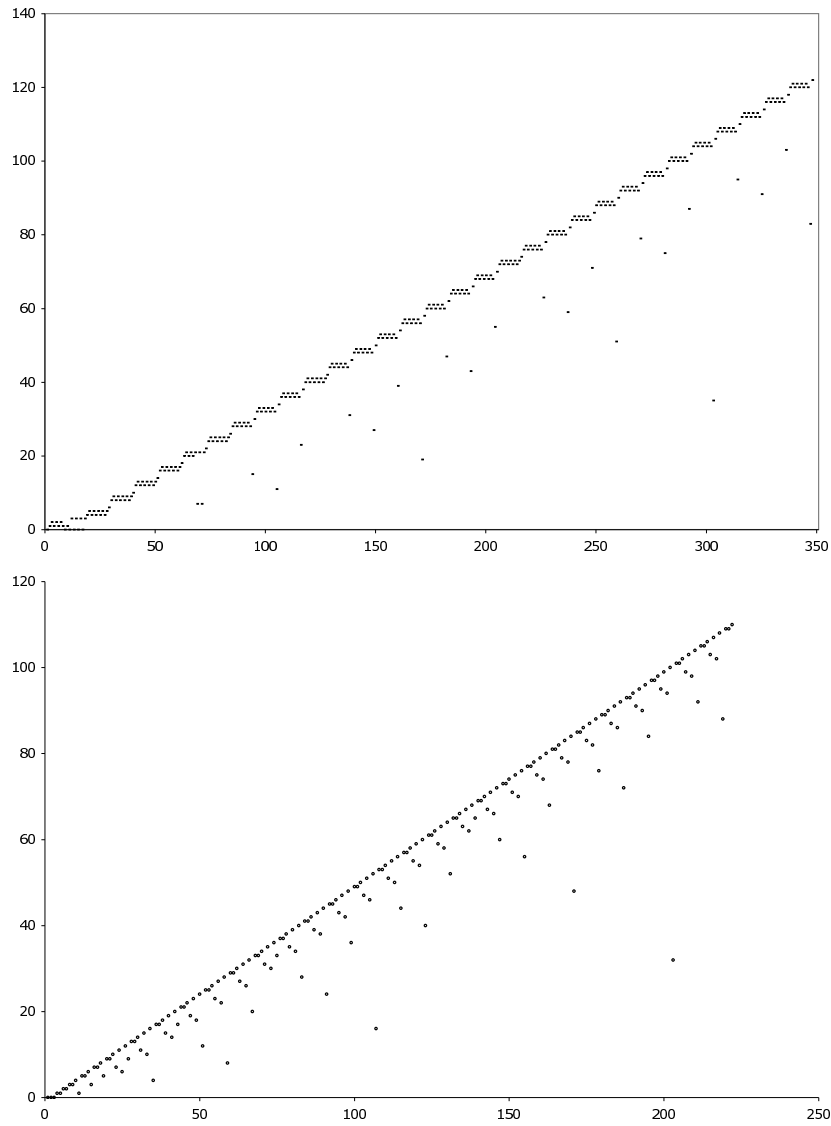
This is illustrated in Figure 1 for **0.2048**, where  $r = 13$ . New “lines” start at  $(m, q) = (1, 0), (2, 0), \dots$ , and are indicated in the top graph. A ruler-regular nim-sequence appears arithmetic-periodic regardless of how far we extend the sequence, but before the sequence goes far enough to ensure arithmetic-periodicity (approximately  $3e + 8ps^3$ , where  $e$  is the size of the largest heap which is not in the period; see [Howse and Nowakowski 2004, Theorem 4]), a new term arises which essentially doubles the length of the apparent period. In this paper we show that **0.20...048**, where there are an odd number of 0s in the hexadecimal code, are ruler-regular. These are not the only such games, for instance **0.21317809532**, **0.31711188** (Figure 2, top), **0.321432132900903213**,



**Figure 1.** The first 350 values (top) and the first 4000 (bottom) of the  $\mathcal{G}$ -sequence of **0.2048**.

and **0.404008** (Figure 2, bottom) all appear to be ruler-regular with  $|K| = 1$ . (Actually, **0.404008** isn't precisely ruler-regular; the sharp-eyed reader will have noticed the same behaviour as in **0.123456789** happening along the first diagonal underneath the main diagonal, in tandem with the point starting the new "line".) The game **0.9138B835B** has  $|K| = 3$ ; see Figure 3.

All of the previously known regularities — periodic, arithmetic-periodic and sapp — have the property that only a finite number of nim-values have to be



**Figure 2.** Top: **0.31711188**. Bottom: **0.404008**.

calculated to identify the type of regularity. However, as yet we have no similar mechanism to check for ruler-regularity.

## 2. The hexadecimal game **0.2048**

In the hexadecimal game **0.2048**, the legal moves are to:

- subtract 1 from a heap of size  $\geq 2$ ;

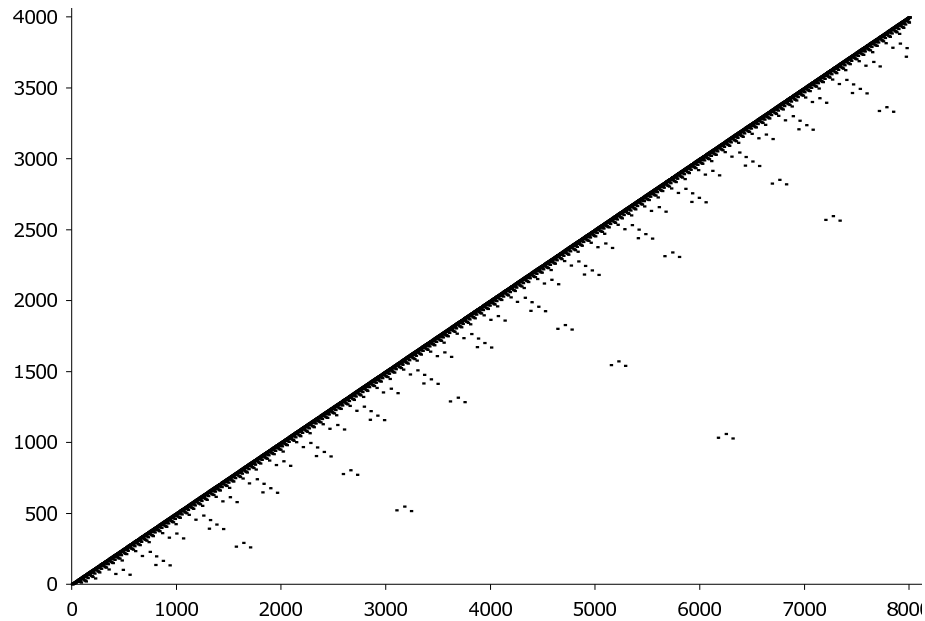


Figure 3. 0.9138B835B.

- subtract 3 from a heap and split the remainder into exactly two heaps;
- subtract 4 from a heap and split the remainder into exactly three heaps.

Here are the first 21 nim-values for this game; they can be computed fairly easily by hand.

$n$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
$\mathcal{G}(n)$	0	0	1	0	1	2	0	1	0	1	2	3	2	3	4	5	4	5	3	4	5

**Theorem 2.** For the game **0.2048**:

- (a) If  $k \geq 0$  and  $j \in \{1, 2, 3, 4, 6, 7, 8, 9, 10, 11, 12, 13, 18\}$ , then  $\mathcal{G}(13k + j) = 4k + \mathcal{G}(j)$  unless  $j = 2$  and  $k$  is of the form

$$(q + 1)2^{m+1} + 2^m + 1, \quad (q, m \geq 0).$$

- (b) For  $q, m \geq 0$ ,  $\mathcal{G}(13((q + 1)2^{m+1} + 2^m + 1) + 2) = 2^{m+3}q + 2^{m+2} + 2$ .

Part (b) of the theorem describes the exceptional values which form the ruler pattern. Equivalently stated, if

$$k = (q + 1)2^{m+1} + 2^m + 1$$

then

$$\mathcal{G}(13k + 2) = 4(k - 2^{m+1} - 1) + 2.$$

$m$	$q = 0$	$q = 1$	$q = 2$	$q = 3$	$q = 4$	Differences
0	(54,6)	(80,14)	(106,22)	(132,30)	(158,38)	(13·2,8)
1	(93,10)	(145,26)	(197,42)	(249,58)	(301,74)	(13·4,16)
2	(171,18)	(275,50)	(379,82)	(483,114)	(587,146)	(13·8,32)
3	(327,34)	(535,98)	(743,162)	(951,226)	(1159,290)	(13·16,64)
4	(639,66)	(1055,194)	(1471,322)	(1887,450)	(2303,578)	(13·32,128)
5	(1263,130)	(2095,386)	(2927,642)	(3759,898)	(4591,1154)	(13·64,256)
6	(2511,258)	(4175,770)	(5839,1282)	(7503,1794)	(9167,2306)	(13·128,512)

**Table 1.** The “lines”  $m = 0, 1, \dots, 6$ .

These exception values form lines, each of slope  $\frac{4}{13}$ , indexed by  $m \geq 0$  starting at  $(39 \times 2^m + 15, 2^{m+2} + 2)$ , and the other points on the “line” are given by  $(39 \times 2^m + 15, 2^{m+2} + 2) + 13(2^{m+1}, 2^{m+3})$ . This is illustrated in Table 1. The annotations “ $m =$ ” in Figure 1 indicate the start of the  $m$ -th line. In the top half of the figure, the four points  $m = 4, q = 0$  through  $m = 1, q = 7$  lie on a “rule” of negative slope. For a given  $k \geq 0$ , starting with the point in the bottom right position of this line, the coordinates of the  $k$ -th new rule are given by Theorem 2(b) with

$$(m, q) = (k, 0), (k - 1, 1), (k - 2, 3), \dots, (k - i, 2^i - 1), \dots, (0, 2^k - 1).$$

The difference between consecutive points is  $(-13 \cdot 2^{m-1}, 2^{m+1})$  which gives a slope of  $-\frac{4}{13}$ . The other set of interesting points are the ones given by  $(m, 0)$ ,  $m = 1, 2, \dots$ . These also fall on a line and this has slope  $\frac{4}{3 \cdot 13}$ .

Before we embark on the proof of Theorem 2, let us take a moment to discuss the statement of the theorem and make a few simple observations that will be of use later on. First, the set  $J = \{1, 2, 3, 4, 6, 7, 8, 9, 10, 11, 12, 13, 18\}$  is a modified set of remainders (mod 13) chosen so that for  $j \in J$  we have  $0 \leq \mathcal{G}(j) \leq 3$ . These are the values of  $\mathcal{G}(j)$  for  $j \in J$ :

$j$	1	2	3	4	6	7	8	9	10	11	12	13	18
$\mathcal{G}(j)$	0	1	0	1	0	1	0	1	2	3	2	3	3

Our proof of Theorem 2 is by induction, where we divide into cases based on  $j$ , the residue modulo 13. For each  $j \in J$ , we must show that  $\mathcal{G}(13k + j)$  is as described by the theorem. In order to show that  $\mathcal{G}(n) = g$ , it is necessary to (i) prove that for all  $g' < g$  there is a move from a heap of size  $n$  to a position with nim-value  $g'$ , and (ii) prove that there is no move from a heap of size  $n$  to a position with nim-value  $g$ . Note that (a) holds trivially with  $k = 0$ , so in all that follows, let  $k > 0$  and  $j \in J$ , and suppose inductively that the statement of Theorem 2 holds for heaps smaller than  $13k + j$ . We begin with a lemma:

**Lemma 3.** *For all  $1 \leq a \leq k$  there is a move from a heap of size  $13k + j$  to a position with nim-value  $4(k - a) + 2$  unless  $k = (q + 1)2^{m+1} + 2^m + 1$ ,  $j = 2$  and  $a = 2^{m+1} + 1$ , in which case there is no move from a heap of size  $13k + j$  to a position with nim-value  $4(k - a) + 2$ .*

*Proof.* We first consider the nonexceptional cases where  $j \neq 2$  or  $a \neq 2^{m+1} + 1$  or  $k \neq (q + 1)2^{m+1} + 2^m + 1$ , and show that there is a move from a heap of size  $13k + j$  to a position with nim-value  $4(k - a) + 2$ . We can subtract 4 and split the remainder into three heaps of size  $r$ ,  $r$  and  $13k + j - 4 - 2r$ . The nim-value of this position is simply  $\mathcal{G}(13k + j - 4 - 2r)$ , so it certainly suffices to find  $r \geq 1$  such that  $\mathcal{G}(13k + j - 4 - 2r) = 4(k - a) + 2$ . In other words, it suffices to find  $n \leq 13k + j - 6$  such that  $13k + j - n \equiv 0 \pmod{2}$  and  $\mathcal{G}(n) = 4(k - a) + 2$ . For  $a = k$ , either  $n = 5$  or  $n = 10$  will satisfy these conditions unless  $13k + j = 14$ , but for this case we simply subtract 3 and divide into heaps of 10 and 1. For  $a < k$  we consider two cases:

**Case 1:  $j + a \equiv 0 \pmod{2}$ .** Take  $n = 13(k - a) + 10$ . Then by induction,  $\mathcal{G}(n) = 4(k - a) + 2$ , and also  $13k + j - n \equiv j + a \equiv 0 \pmod{2}$ , so we are done if  $13(k - a) + 10 \leq 13k + j - 6$ , which is equivalent to  $16 \leq 13a + j$ . This is certainly the case if  $a > 1$ , so we need only to consider  $a = 1$  in which case the inequality becomes  $j \geq 3$ . This in turn is true unless  $j = 1$  or  $j = 2$ , but  $j \neq 2$  because we are considering the case  $j + a \equiv 0 \pmod{2}$ . Thus, it remains to show that there is a move from a heap of size  $13k + 1$  to a position with nim-value  $4(k - 1) + 2$ . We can accomplish this by subtracting 3 and splitting the remainder into two heaps of size 1 and  $13k - 3$ ; by induction  $\mathcal{G}(1) \oplus \mathcal{G}(13k - 3) = 0 \oplus \mathcal{G}(13(k - 1) + 10) = 4(k - 1) + 2$ .

**Case 2:  $j + a \equiv 1 \pmod{2}$ .** Let  $m$  be the number of trailing zeros in the binary representation of  $k - a$ . Theorem 2(b) states that  $4(k - a) + 2$  appears as an exceptional value  $\mathcal{G}(n)$  for  $n = 13(k - a + 2^{m+1} + 1) + 2$ . If either  $a > 2^{m+1} + 1$  or  $a = 2^{m+1} + 1$  and  $j > 2$  then  $n = 13(k - a + 2^{m+1} + 1) + 2 < 13k + j$ , so by induction we can assume that  $\mathcal{G}(n) = 4(k - a) + 2$ . Then there is some  $q$  such that  $4(k - a) + 2 = 2^{m+3}q + 2^{m+2} + 2$ , and so  $13k + j - n \equiv j + a - 1 \equiv 0 \pmod{2}$ , so we're done if  $n \leq 13k + j - 6$ , that is,  $8 \leq 13(a - 2^{m+1} - 1) + j$ . This in turn is true if either  $a > 2^{m+1} + 1$  or  $a = 2^{m+1} + 1$  and  $j \geq 8$ ; note that these conditions on  $(a, j)$  subsume the previous ones. Thus, we may assume from this point forward that  $a \leq 2^{m+1} + 1$  and either  $a < 2^{m+1} + 1$  or  $j < 8$ . In particular, for  $j \geq 8$  we may assume that  $a < 2^{m+1} + 1$ . Note that by definition of  $m$ , if  $d < 2^m$  then the 1s in the binary representations of  $k - a$  and  $d$  do not overlap, so  $d \oplus (k - a) = d + (k - a)$ , which we rewrite as

$$k - a = d \oplus (k - a + d). \quad (1)$$



We proceed by considering three cases.

**Case 2.1:  $a \equiv 0 \pmod{2}$ ,  $j \equiv 1 \pmod{2}$ .** Since  $a \geq 1$  we in fact have  $a \geq 2$ . Let  $d = (a - 2)/2$ . Then  $d \leq (2^{m+1} - 1)/2 < 2^m$  and (1) becomes

$$k - a = d \oplus (k - 2 - d). \quad (2)$$

Our strategy is now to subtract 4 and split the remaining  $13k + j - 4$  into exactly three heaps of sizes  $13(k - 2 - d) + j_1$ ,  $13d + j_2$  and  $j_3$  where  $j_1 + j_2 + j_3 = 22 + j$ ,  $\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = 2$ ,  $j_3 \in J \cup \{5\}$  and  $j_1, j_2 \in J$ . Using induction, (2) and the fact that  $0 \leq \mathcal{G}(j_i) \leq 3$ , the nim-value of this position is

$$\begin{aligned} & (4(k - 2 - d) + \mathcal{G}(j_1)) \oplus (4d + \mathcal{G}(j_2)) \oplus \mathcal{G}(j_3) \\ &= 4(d \oplus (k - 2 - d)) + \mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = 4(k - a) + 2. \end{aligned}$$

Since  $j \equiv 1 \pmod{2}$ , we only need to consider  $j \in \{1, 3, 7, 9, 11, 13\}$ . We can verify that the  $j_i$  in the following table satisfy the given conditions for each choice of  $j$ :

$22 + j$	23	25	29	31	33	35
$j_1$	9	10	12	13	7	8
$j_2$	9	10	12	13	8	9
$j_3$	5	5	5	5	18	18

**Case 2.2a:  $a \equiv 1 \pmod{2}$ ,  $j \equiv 0 \pmod{2}$ ,  $a < 2^{m+1} + 1$ ,  $j \geq 8$  or  $j = 2$ .** Let  $d = (a - 1)/2 < 2^{m+1}/2 = 2^m$ . Then (1) becomes

$$k - a = d \oplus (k - 1 - d). \quad (3)$$

The conditions on  $j$  imply that  $j \in \{2, 8, 10, 12, 18\}$  and hence

$$(j + 4)/2 \in \{3, 6, 7, 8, 11\} \subset J.$$

We can therefore subtract 4 and split the remaining  $13k + j - 4$  into three heaps of size  $13(k - 1 - d) + (j + 4)/2$ ,  $13d + (j + 4)/2$  and 5. Using induction and (3), we again find that the nim-value of this position is  $4(k - a) + 2$ .

**Case 2.2b:  $a \equiv 1 \pmod{2}$ ,  $a \leq 2^{m+1} + 1$ ,  $j = 4, 6$ .** We dispense with the case  $a = 1$  by observing that we can subtract 4 from  $13k + 4$  and split the remaining  $13k$  into heaps of size 5, 4 and  $13(k - 1) + 4$ ; we can subtract 3 from  $13k + 6$  and split the remaining  $13k + 3$  into heaps of size 6 and  $13(k - 1) + 10$ . By induction, the nim-value of each resulting position is  $4(k - 1) + 2$ . Now assume that  $a \geq 3$ , and let  $d = (a - 3)/2 \leq (2^{m+1} - 2)/2 < 2^m$ . Then (1) becomes

$$k - a = d \oplus (k - 3 - d). \quad (4)$$

Our strategy is to again subtract 4 and split the remaining  $13k + j - 4$  into exactly three heaps of size  $13(k - 3 - d) + j_1$ ,  $13d + j_2$  and  $j_3$  where  $j_1 + j_2 + j_3 = 35 + j$ ,  $\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = 2$ ,  $j_3 \in J \cup \{5\}$  and  $j_1, j_2 \in J$ . Using induction, (4) and the fact that  $0 \leq \mathcal{G}(j_i) \leq 3$ , we again find that the nim-value of this position is  $4(k - a) + 2$ . We can verify that for  $j = 4$  we can take  $(j_1, j_2, j_3) = (10, 11, 18)$  and for  $j = 6$  we can take  $(j_1, j_2, j_3) = (11, 12, 18)$ .

The only remaining values of  $j$  and  $a$  to consider are  $j = 2$  and  $a = 2^{m+1} + 1$ ; in all other cases we have shown that there is a move from a heap of size  $13k + j$  to a position with nim-value  $4(k - a) + 2$ . Recall that  $m$  is the number of trailing zeros in the binary representation of  $k - a$ , thus  $k - a$  is of the form  $q \cdot 2^{m+1} + 2^m$  for some  $q \geq 0$ , that is  $k = (k - a) + a = (q + 1) \cdot 2^{m+1} + 2^m + 1$ . This is exactly the case excepted by the statement of the lemma, so all that remains to be shown is that in this case there is no move from a heap of size  $13k + j$  to a position with nim-value  $4(k - a) + 2$ . First, observe that subtracting 1 leaves a heap which by induction has nim-value  $4k$ . Next, suppose we subtract 4 and divide the remaining  $13k - 2$  into exactly 3 heaps. We can express the size of these heaps as  $13k_1 + j_1$ ,  $13k_2 + j_2$ , and  $13k_3 + j_3$ , where for  $i = 1, 2, 3$  either  $j_i \in J$  or  $j_i = 5$  and  $k_i = 0$ . By induction, the nim-value of the resulting position is:

$$\begin{aligned} & (4k_1 + \mathcal{G}(j_1)) \oplus (4k_2 + \mathcal{G}(j_2)) \oplus (4k_3 + \mathcal{G}(j_3)) \\ & = 4(k_1 \oplus k_2 \oplus k_3) + \mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3). \end{aligned}$$

If this is equal to  $4(k - a) + 2$  then the following three equations must hold:

$$(13k_1 + j_1) + (13k_2 + j_2) + (13k_3 + j_3) = 13k - 2, \quad (5)$$

$$k_1 \oplus k_2 \oplus k_3 = k - a, \quad (6)$$

$$\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = 2. \quad (7)$$

From (5) we have  $j_1 + j_2 + j_3 \equiv -2 \pmod{13}$ . On the other hand, since  $j_i \leq 18$  we have  $j_1 + j_2 + j_3 \leq 54$ , thus  $j_1 + j_2 + j_3 \in \{11, 24, 37, 50\}$ . There are no choices of  $j_1, j_2, j_3 \in J$  which sum to 50. If  $j_1 + j_2 + j_3 = 37$  then they can't all be  $\leq 13$ , so one of them is 18 and the other two sum to 19, but we can verify by hand that for all such choices we have  $\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = 0$ . If  $j_1 + j_2 + j_3 = 24$  then from (5)  $k_1 + k_2 + k_3 = k - 2$ . But  $a = 2^{m+1} + 1$  is odd and  $k_1 \oplus k_2 \oplus k_3 = k - a$ , so in this case  $k_1 + k_2 + k_3$  and  $k_1 \oplus k_2 \oplus k_3$  would have opposite parity which is impossible. Finally, if  $j_1 + j_2 + j_3 = 11$  then each  $j_i \leq 9$ , so in order to satisfy (7) one of them must be 5, say  $j_3$ , hence  $k_3 = 0$ . We are then left with two equations:

$$k_1 + k_2 = k - 1, \quad (8)$$

$$k_1 \oplus k_2 = k - a. \quad (9)$$

Now  $k_1 + k_2 = k_1 \oplus k_2 + 2(k_1 \& k_2)$ , where  $\&$  denotes bitwise Boolean AND. Equations (8) and (9) therefore give us  $2(k_1 \& k_2) = k_1 + k_2 - k_1 \oplus k_2 = a - 1 = 2^{m+1}$  which implies that  $k_1 \& k_2 = 2^m$  and so  $k_1$  and  $k_2$  both have a 1 in the  $2^m$  place of their binary representations. But  $k_1 \oplus k_2 = k - a = q \cdot 2^{m+1} + 2^m$  so that exactly one of  $k_1$  and  $k_2$  have a 1 in the  $2^m$  place of their binary representations, a contradiction. Hence there is no way to subtract 4 and divide the remaining  $13k - 2$  into three heaps such that the resulting position has nim-value  $4(k - a) + 2$ . Finally, suppose we subtract 3 and divide the remaining  $13k - 1$  into exactly 2 heaps of size  $13k_1 + j_1$  and  $13k_2 + j_2$  where again either  $j_i \in J$  or  $j_i = 5$  and  $k_i = 0$ . By induction, the nim-value of the resulting position is  $4(k_1 \oplus k_2) + \mathcal{G}(j_1) \oplus \mathcal{G}(j_2)$ . We therefore have the following three equations, corresponding to Equations (5)–(7):

$$(13k_1 + j_1) + (13k_2 + j_2) = 13k - 1, \quad (10)$$

$$k_1 \oplus k_2 = k - a, \quad (11)$$

$$\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) = 2. \quad (12)$$

From (10) we have  $j_1 + j_2 \equiv -1 \pmod{13}$ , but  $j_i \leq 18$  so that  $j_1 + j_2 \leq 36$  giving  $j_1 + j_2 \in \{12, 25\}$ . Hence from (10),  $k_1 + k_2 = k - 1$  or  $k_1 + k_2 = k - 2$ . But we saw previously that  $k_1 + k_2 = k - 1$  is inconsistent with (11) (which is the same as (9)), and  $k - 2$ ,  $k - a$  have opposite parity. Hence there is no way to subtract 3 and divide the remaining  $13k - 1$  into two heaps such that the resulting position has nim-value  $4(k - a) + 2$ , and this concludes the proof of Lemma 3.  $\square$

**Lemma 4.** *For all  $a \geq 1$ , there is a move from a heap of size  $13k + j$  to positions with nim-values  $4(k - a) + 1$ ,  $4(k - a)$  and  $4(k - a) - 1$ .*

*Proof.* As in the proof of the previous lemma, we can move to a position with a given nim-value if a heap of size  $n$  has the desired value where  $n \leq 13k + j - 6$  and  $13k + j - n \equiv 0 \pmod{2}$ . Since  $j \geq 1$ , this will certainly be the case if  $n \leq 13(k - 1) + 8$  and  $13k + j - n$  is even. Hence, we obtain nim-value  $4(k - a) + 1$  with either  $n = 13(k - a) + 7$  or  $n = 13(k - a) + 4$  (one or the other will always do since they have opposite parity), we obtain nim-value  $4(k - a)$  with either  $n = 13(k - a) + 6$  or  $n = 13(k - a) + 3$ , and we obtain nim-value  $4(k - a) - 1$  with either  $n = 13(k - a - 1) + 11$  or  $n = 13(k - a - 1) + 18$ .  $\square$

We now know that from a heap of size  $13k + j$ , we can move to a position with nim-value  $g$  for any  $g \leq 4k - 2$ , with one exception given by Lemma 3. This in fact proves part (b) of Theorem 2, as it shows that from a heap of size

$$13((q + 1) \cdot 2^{m+1} + 2^m + 1) + 2,$$

we cannot move to a position with value  $2^{m+3}q + 2^{m+2} + 2$ , but we can move to a position with value  $g$  for all  $g < 2^{m+3}q + 2^{m+2} + 2$ , so

$$\mathcal{G}(13((q+1) \cdot 2^{m+1} + 2^m + 1) + 2) = 2^{m+3}q + 2^{m+2} + 2.$$

To finish the proof of Theorem 2, we must show that for all nonexceptional heaps  $13k + j$  we can move to a position with value  $g$  for  $4k - 1 \leq g < 4k + \mathcal{G}(j)$ , but we cannot move to a position with value  $4k + \mathcal{G}(j)$ . Table 2 accomplishes the former, giving explicit moves for each  $j \in J$ .

It remains to show that there is no move from a heap of size  $13k + j$  to a position with nim-value  $4k + \mathcal{G}(j)$ . First, subtracting 1 leaves  $13k + j - 1$ . By induction, if  $j - 1 \in J$ , then  $\mathcal{G}(13k + j - 1) = 4k + \mathcal{G}(j - 1)$ , but  $\mathcal{G}(j) \neq \mathcal{G}(j - 1)$  for  $j, j - 1 \in J$ ; otherwise  $j + 12 \in J$  and

$$\mathcal{G}(13k + j - 1) = \mathcal{G}(13(k - 1) + j + 12) = 4(k - 1) + \mathcal{G}(j + 12) < 4k + \mathcal{G}(j).$$

This leaves us with two cases to consider: moves that split  $13k + j - 4$  into exactly three heaps, and moves that split  $13k + j - 3$  into exactly two heaps. Following the approach in the proof of Lemma 3, splitting  $13k + j - 4$  into heaps of size  $13k_i + j_i$  with either  $j_i \in J$  or  $j_i = 5$  and  $k_i = 0$  gives us three equations, for which we must show there is no solution:

$$(13k_1 + j_1) + (13k_2 + j_2) + (13k_3 + j_3) = 13k + j - 4, \quad (13)$$

$$k_1 \oplus k_2 \oplus k_3 = k, \quad (14)$$

$$\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) \oplus \mathcal{G}(j_3) = \mathcal{G}(j). \quad (15)$$

Equation (14) tells us that  $k_1 + k_2 + k_3 \geq k_1 \oplus k_2 \oplus k_3 = k$ , so we can rewrite (13) as:

$$j - 4 - j_1 - j_2 - j_3 = 13(k_1 + k_2 + k_3 - k) \geq 0. \quad (16)$$

For each  $j$  there are limited choices for  $j_1, j_2, j_3$  satisfying both

$$j_1 + j_2 + j_3 \equiv j - 4 \pmod{13} \quad \text{and} \quad j_1 + j_2 + j_3 \leq j - 4,$$

and we can verify by hand that none of them also satisfy (15). The same approach can be used for moves which split  $13k + j - 3$  into two heaps; in this case we verify by hand that there are no  $j_1, j_2$  satisfying  $j_1 + j_2 \equiv j - 3 \pmod{13}$ ,  $j_1 + j_2 \leq j - 3$ , and  $\mathcal{G}(j_1) \oplus \mathcal{G}(j_2) = \mathcal{G}(j)$ . This completes the proof of Theorem 2.

### 3. The games 0.200048, 0.20000048, 0.2000000048, . . .

**Theorem 5.** *Let  $G$  be the game  $0,2(0)^{2n+1}48$  (i.e.,  $2n + 1$  zeroes between the 2 and 4) and let  $A = 6(n + 2) + 1$ .*

$j$	Move	Nim-value
1	$(13k + 1) - 1 = (13(k - 1) + 13)$	$4(k - 1) + 3 = 4k - 1$
2	$(13k + 2) - 3 = (1) + (13(k - 1) + 11)$ $(13k + 2) - 1 = (13k + 1)$	$4(k - 1) + 3 = 4k - 1$ $4k$
3	$(13k + 3) - 4 = (3) + (5) + (13(k - 1) + 4)$	$4(k - 1) + 3 = 4k - 1$
4	$(13k + 4) - 4 = (1) + (1) + (13(k - 1) + 11)$ $(13k + 4) - 1 = (13k + 3)$	$4(k - 1) + 3 = 4k - 1$ $4k$
6	$(13k + 6) - 1 = (13(k - 1) + 18)$	$4(k - 1) + 3 = 4k - 1$
7	$(13k + 7) - 3 = (6) + (13(k - 1) + 11)$ $(13k + 7) - 1 = (13k + 6)$	$4(k - 1) + 3 = 4k - 1$ $4k$
8	$(13k + 8) - 4 = (2) + (2) + (13(k - 1) + 13)$	$4(k - 1) + 3 = 4k - 1$
9	$(13k + 9) - 3 = (8) + (13(k - 1) + 11)$ $(13k + 9) - 1 = (13k + 8)$	$4(k - 1) + 3 = 4k - 1$ $4k$
10	$(13k + 10) - 4 = (3) + (3) + (13(k - 1) + 13)$ $(13k + 10) - 3 = (1) + (13k + 6)$ $(13k + 10) - 1 = (13k + 9)$	$4(k - 1) + 3 = 4k - 1$ $4k$ $4k + 1$
11	$(13k + 11) - 3 = (3) + (13(k - 1) + 18)$ $(13k + 11) - 4 = (2) + (2) + (13k + 3)$ $(13k + 11) - 3 = (1) + (13k + 7)$ $(13k + 11) - 1 = (13k + 10)$	$4k - 1$ $4k$ $4k + 1$ $4(k - 1) + 3 = 4k + 2$
12	$(13k + 12) - 4 = (4) + (4) + (13(k - 1) + 13)$ $(13k + 12) - 4 = (1) + (1) + (13k + 6)$ $(13k + 12) - 4 = (2) + (2) + (13k + 4)$	$4(k - 1) + 3 = 4k - 1$ $4k$ $4k + 1$
13	$(13k + 13) - 4 = (1) + (3) + (13(k - 1) + 18)$ $(13k + 13) - 4 = (3) + (3) + (13k + 3)$ $(13k + 13) - 3 = (3) + (13k + 7)$ $(13k + 13) - 1 = (13k + 12)$	$4(k - 1) + 3 = 4k - 1$ $4k$ $4k + 1$ $4k + 2$
18	$(13k + 18) - 4 = (7) + (7) + (13(k - 1) + 13)$ $(13k + 18) - 4 = (3) + (3) + (13k + 8)$ $(13k + 18) - 4 = (5) + (5) + (13k + 4)$ $(13k + 18) - 4 = (1) + (1) + (13k + 12)$	$4(k - 1) + 3 = 4k - 1$ $4k$ $4k + 1$ $4k + 2$

**Table 2.** Moving to a position with nim-value  $g$ ,  $4k - 1 \leq g < 4k + \mathcal{G}(j)$ .

- (a) If  $k > 0$  then  $\mathcal{G}(Ak + j) = 4k + \mathcal{G}(j)$  unless  $j = 2$  and  $k$  is of the form  $(q + 1) \cdot 2^{m+1} + 2^m + 1$ ,  $q, m \geq 0$ .
- (b) For  $q, m \geq 0$ ,  $\mathcal{G}(A((q + 1) \cdot 2^{m+1} + 2^m + 1) + 2) = 2^{m+3}q + 2^{m+2} + 2$ .

*Proof.* It is not too difficult to prove that the generic game  $\mathbf{0.2(0)^{2n+1}48}$  has initial nim-sequence  $0, (0, 1)^{n+2}, 2, (0, 1)^{n+2}, (2, 3)^{n+2}$ . The proof is now a repetition of that for  $\mathbf{0.2048}$ , with fewer special cases to be considered. We leave this to the reader.  $\square$

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